

# COSC 2011 Section N

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## Overview

- Mathematical Induction
  - ◆ Definition, Examples
- Introduction to Recursion
  - ◆ Definition, Examples
- Assignment 1 Notes/Questions?
  - ◆ For loop time analysis

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## Loop Invariants: (1)

- An assertion that remains true each time the statements of a loop are executed
  - ◆ Tells us something about the values of the loop variables while executing a loop.
  - ◆ Should be true at the beginning of each iteration, including the first!

$\{p\} \text{ while } C \text{ do } S \rightarrow \{p \wedge \neg C\}$

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## Bubble Sort Algorithm (1)

**Algorithm** Bubblesort(*sequence*):

*Input:* sequence of integers *sequence*

*Postcondition:* *sequence* is sorted & contains the same integers as the original sequence

*length* = length of *sequence*

```
for i = 0 to length - 1 do
  for j = 0 to length - i - 2 do
    if jth element of sequence >
      (j+1)th element of sequence
    then
      swap jth and (j+1)th element
      of sequence
```

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## Bubble Sort Algorithm (2)

- Loop Invariant – Outer Loop:
  - ◆ Last *i* elements of *sequence* are sorted and are all greater or equal to the other elements of the sequence.
- Loop Invariant – Inner Loop:
  - ◆ Same as outer loop and the *j*th element of *sequence* is greater or equal to the first *j* elements of *sequence*.

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### Bubble Sort Algorithm (3)

- Running Time Analysis:
    - ◆ Assume access to and swap of elements takes  $O(1)$  time.
    - ◆ Running time of  $i$ th pass:  
 $O(\text{sum}[n-i+1])$
    - ◆ Can re-write it as:  
 $O(n + (n+1) + \dots + 2 + 1)$   
 $O(\text{sum}(i))$
- By proposition 3.4:  
 $\text{sum}(i) = [n(n+1)] / 2$

### Mathematical Induction: (1)

- Can be used only to prove results obtained some other way:
  - ◆ Not a tool for discovering formulas or theorems!
- Many theorems state  $P(n)$  is true for all positive integers  $n$ 
  - ◆ Mathematical induction is used to prove assertions (propositions) of this kind.

### Mathematical Induction: (2)

- Used to prove statements of the form  $\forall n P(n)$ , for all positive integers.
- A proof by mathematical induction that  $P(n)$  is true for all positive integers consists of two steps
  1. Basis Step: show  $P(1)$  (or  $n = \text{some other finite value}$ ) is true.
  2. Inductive Step: Show  $P(n) \rightarrow P(n+1)$  is true for every positive integer  $n$ .

### Mathematical Induction: (3)

- $P(n)$  is called the *inductive hypothesis*. When both steps are done, then we have shown  $\forall n P(n)$ .

$$[P(1) \wedge \forall n (P(n) \rightarrow P(n+1))] \rightarrow \forall n P(n)$$

- ◆ To prove the inductive step for every  $n$ , we need to show  $P(n)$  cannot be false when  $P(n)$  is true.
  - ★ Assume  $P(n)$  is true & show that under this assumption,  $P(n+1)$  must be true.

### Mathematical Induction: (4)

- **Remark:** It is not assumed  $P(n)$  is true for all positive integers! Only shown that if it is assumed  $P(n)$  is true then  $P(n+1)$  is also true.
- When using induction, we show that  $P(1)$  is true. Then since  $P(1)$  implies  $P(2)$ ,  $P(2)$  must be true. Then  $P(3)$  is true because  $P(2)$  implies  $P(3)$ . Continuing along these lines,  $P(k)$  is true for any positive integer  $k$ .

### Mathematical Induction: (5)

- **Useful Illustration:**
  - ◆ Consider a line of people, person 1, person 2 etc. A secret is told to the first person and each person tells the secret to the next person in line.
  - ◆ Let  $P(n)$  be the statement that person  $n$  knows the secret.
    - ★  $P(1)$  is true since it was told to first person.
    - ★  $P(2)$  is true since person 1 tells person 2 and so on...

### Mathematical Induction: (6)

- **Another Illustration:**
  - ◆ Infinite row of dominos, labeled  $1, 2, 3, \dots, n$  & each domino is standing up.
  - ◆ Let  $P(n)$  be the statement that domino  $n$  is knocked over.
  - ◆ If the first domino is knocked over,  $P(1)$  is true.
  - ◆ If whenever first domino is knocked over –  $P(1)$  is true, it knocks the  $(n+1)$ th domino over –  $P(n) \rightarrow P(n+1)$  is true, then all dominos are knocked over!

### Mathematical Induction: (7)

- Sometimes we need to show  $P(n)$  is true for  $n=k, k+1, k+2, \dots$  where  $k$  is an integer other than 1.
  - ◆ Can still use induction as long as we change the *Basis Step*.
  - ◆ Show  $P(k)$  is true and then show  $P(n) \rightarrow P(n+1)$  is true for  $n = k, k+1, k+2, \dots$
  - ◆  $k$  can be negative, positive or zero.

### Recursive Methods (1):

- Sometimes we can reduce the solution to a problem with a particular input to the solution of the same problem with smaller input.
- ◆ Solution to the original problem can be found with a sequence of **reductions** until problem is reduced to some initial case where solution is known.

### Recursion (Math): (2):

- Definition (Mathematical):
  - ◆ When an object is defined in terms of itself.
- Can be Used to Define:
  - ◆ Sequences, functions sets.
- Example:
  - ◆ Sequence of powers of 2 is given by  $a_n = 2^n$
  - ◆ Can also be defined as:
    - ★ Give the first term of the sequence:  $a_0 = 1$

### Recursion (Math): (3):

- ★ Rule for finding a term of the sequence from the previous one.
- Recursively Defined Functions:
  1. Specify the value of the function at 0.
  2. Give rule for finding its value as an integer from its values at smaller integers.

### Recursion (Math): (4):

- Example:
  - ◆  $f$  is defined recursively as:  
 $f(0) = 3$   
 $f(n + 1) = 2f(n) + 3$   
 $f(1) = 2f(0) + 3 = (2 \times 3) + 3 = 9$   
 $f(2) = 2f(1) + 3 = (2 \times 9) + 3 = 21$   
 $f(3) = 2f(2) + 3 = (2 \times 21) + 3 = 45$
- Question:
  - ◆ Give an inductive definition of the factorial function:  $f(n) = n!$

### Recursion (Math): (5):

- Solution:
  - ◆ Initial value:  
 $f(0) = 1$
  - ◆ Rule for finding  $f(n+1)$ :
    - ★  $(n+1)!$  is computed by multiplying  $n!$  by  $(n+1)$   
 $f(n+1) = f(n) \times (n+1)$
- Determining the Value of the Factorial Function:
  - ◆ Use the rule that shows how to express  $f(n+1)$

### Recursion (Math): (6):

- ◆ In terms of  $f(n)$  several times:
  - Example of  $f(4) = 4!$   
 $f(4) = 4f(3) = 4 \times 3f(2) = 4 \times 3 \times 2f(1) = 4 \times 3 \times 2 \times 1 \times f(0) = 4 \times 3 \times 2 \times 1 \times 1 = 24$
  - ◆ When  $f(0)$  is the only function that occurs, no more reductions are necessary.
    - ★ Only thing to do is insert  $f(0)$  into formula.

### Recursion (Math): (7):

- Example: Fibonacci Numbers  $f_0, f_1, f_2, \dots, f_n$  are defined as follows:  
 $f_0 = 0,$   
 $f_1 = 1$  and  
 $f_n = f_{n-1} + f_{n-2} \quad n = 2, 3, 4, \dots$
- ◆ What are  $f_2, f_3, f_4, f_5$ ?

### Recursion (Math): (8):

- ◆ Solution to Fibonacci Numbers for  $n = 2, 3, 4, 5, 6$   
 $f_2 = f_1 + f_0 = 1 + 0 = 1$   
 $f_3 = f_2 + f_1 = 1 + 1 = 2$   
 $f_4 = f_3 + f_2 = 2 + 1 = 3$   
 $f_5 = f_4 + f_3 = 3 + 2 = 5$   
 $f_6 = f_5 + f_4 = 5 + 3 = 8$

### Recursive Methods (1):

- A method is called *recursive* if it solves a problem by reducing it to an instance of the same problem with smaller input.
- ◆ A Recursive Method calls itself as a subroutine with.
- ◆ Each time it calls itself, the problem is “reduced” until we reach a point where the problem is small enough to be easily solved.

### Recursive Methods (2):

- Candidate Problems for Recursion Have the Following Characteristics:
  - ◆ One or more simple cases of the problem have simple non-recursive solution (*base cases*).
  - ◆ For other cases there is a process for substituting one or more reduced cases of the problem that are closer to the base case.

### Recursive Methods (3):

- ◆ Eventually the problem can be reduced to base cases only, all of which are easy to solve!
- ◆ Recursive algorithms we will encounter will generally be of the form:

**If** *base case* reached **then**  
    solve it  
**else**  
    reduce problem using recursion

### Recursive Methods (4):

- Recursive Factorial:

```
public static long factorial(long n){  
    if (n <= 1)  
        return 1;  
    else  
        return n * factorial(n - 1);  
}
```
- ◆ Method calls itself recursively to compute factorial of n-1.
- ◆ When recursive call terminates, returns (n-1)!

### Recursive Methods (5):

- ◆  $(n-1)!$  Is then multiplied by  $n$  to compute  $n!$ .
- ◆ In turn, recursive invocation calls itself to compute the factorial of  $n-2$ , etc...
- ◆ To compute  $n!$ , multiply  $n$  by factorial of  $n-1$ . But how do we calculate factorial( $n-1$ )?
  - ★ We call the factorial method with  $n-1$  as an argument...

### Recursive Methods (6):

- Important Properties Every Recursive Methods Should Possess:
  - ◆ Method must terminate!
    - ★ Base case
    - ★ Even infinite recursive method will terminate!
  - ◆ Always perform the recursive call on a smaller input value.
    - ★ Reduce the problem at every recursive call.

### Recursive Methods (7):

- Example: Fibonacci Numbers

```
public static long fibonacci(long n){
    if (n == 0)
        return 0;
    else if (n == 1)
        return 1;
    else
        return fibonacci(n - 1) +
            fibonacci(n-2);
}
```

### Recursion & Induction:

- Notice the Similarity Between Recursion and Induction!
  - ◆ Induction can be used to prove the correctness of many recursive formulas & functions!
- Problem with Recursion:
  - ◆ Usually require more computation and space over an iterative approach