

COSC3431.03

# TAKE HOME EXAM

Dept. of Computer Science

**Date:** July 13, 1999

**Due:** July 20, 1999, In class, by 7:15pm

- 1. 5 MARKS** Let  $\text{NOPREFIX}(L)$  stand for the set  $\{w \in L : \text{no proper prefix of } w \text{ is a member of } L\}$ .  
Prove that if  $L$  is regular, then  $\text{NOPREFIX}(L)$  is also regular.
- 2. 5 MARKS** We have seen a number of “practical examples” of converting an NFA of  $n$  states to a DFA that “beat” by a substantial margin the “theoretical” upper bound ( $2^n$ ) for the number of states of the DFA. This was due to the fact that many subset-states introduced were *inessential*, i.e., we could remove them *without affecting the accepted language*.  
Nevertheless, show that for any  $n$ , an NFA of  $n$  states exists whose equivalent DFA—constructed by the “subset construction”—has no less than  $c2^n$  *essential* (i.e., non-removable) states, where  $c$  is a small (possibly fractional) constant.
- 3. 5 MARKS** Prove that the class (nondeterministic) PDAs that accept by “simultaneous empty stack and final state”, ESFS-PDA, are exactly as powerful as ES-PDA or FS-PDA.  
Note that the term in quotes is meant to mean that an ESFS-PDA accepts  $x$  iff  $(q_0, x, \$) \vdash^* (q, \varepsilon, \varepsilon)$  where  $q$  is a *final state*.
- 4. 5 MARKS** Prove that the set  $\{x : \text{ran}(\phi_x) \text{ is infinite}\}$  is not semi-recursive.  
*Hint.* See how you can adapt the proof in class, that  $\{x : \phi_x \text{ is a constant}\}$  is not semi-recursive.
- 5. 5 MARKS** Prove that if  $P(y, \vec{x})$  is semi-recursive, then so is  $(\exists y)_{<z} P(y, \vec{x})$ .
- 6. 5 MARKS** Prove that  $Q(\vec{x})$  is semi-recursive iff (note the “iff”!), for some function  $f \in \mathcal{P}$  of the same number of arguments as  $Q$ , it is true that  $Q(\vec{x}) \equiv f(\vec{x}) = 0$  for all  $\vec{x}$ .