York University

Faculty of Science and Engineering

MATH 1090: Facts-List for the December 2010 Examination (to be held Dec 18, 2010)

The following are the axioms of Propositional Calculus: In what follows, A, B, C stand for arbitrary formulae.

Properties of $\equiv$

Associativity of $\equiv$ $\quad ((A \equiv B) \equiv C) \equiv (A \equiv (B \equiv C)) \quad (1)$

Symmetry of $\equiv$ $\quad (A \equiv B) \equiv (B \equiv A) \quad (2)$

Properties of $\bot, \top$

$\top$ vs. $\bot$ $\quad \top \equiv \bot \equiv \bot \quad (3)$

Properties of $\neg$

Introduction of $\neg$ $\quad \neg A \equiv A \equiv \bot \quad (4)$

Properties of $\vee$

Associativity of $\vee$ $\quad (A \vee B) \vee C \equiv A \vee (B \vee C) \quad (5)$

Symmetry of $\vee$ $\quad A \vee B \equiv B \vee A \quad (6)$

Idempotency of $\vee$ $\quad A \vee A \equiv A \quad (7)$

Distributivity of $\vee$ over $\equiv$ $\quad A \vee (B \equiv C) \equiv A \vee B \equiv A \vee C \quad (8)$

Excluded Middle $\quad A \vee \neg A \quad (9)$

Properties of $\wedge$

Golden Rule $\quad A \wedge B \equiv A \equiv B \equiv A \vee B \quad (10)$

Properties of $\rightarrow$

Implication $\quad A \rightarrow B \equiv A \vee B \equiv B \quad (11)$

The **Primary** Boolean rules are:

$$\frac{A, A \equiv B}{B} \qquad (Eqn)$$

and

$$\frac{A \equiv B}{C[\mathbf{p} := A] \equiv C[\mathbf{p} := B]} \qquad (Leib)$$

The following are the Predicate Calculus Axioms:

Any partial generalisation of any formula in groups* Ax1–Ax6 *is an axiom for Predicate Calculus.

Groups **Ax1–Ax6** contain the following schemata:

Ax1. Every tautology.

Ax2. $(\forall \mathbf{x})A \to A[\mathbf{x} := t]$, for any term t.

Ax3. $A \to (\forall \mathbf{x})A$, provided $\mathbf{x}$ is not free in A.

Ax4. $(\forall \mathbf{x})(A \to B) \to (\forall \mathbf{x})A \to (\forall \mathbf{x})B$.

Ax5. For *each* object variable $\mathbf{x}$, the formula $\mathbf{x} = \mathbf{x}$.

Ax6. For any terms t, s, the schema $t = s \to (A[\mathbf{x} := t] \equiv A[\mathbf{x} := s])$.

The following metatheorems are good for ***both*** Propositional and Predicate Calculus:

1. *Redundant True.* $\Gamma \vdash A$ iff $\Gamma \vdash A \equiv \top$
2. *Modus Ponens (MP).* $A, A \to B \vdash B$
3. *Cut Rule.* $A \vee B, \neg A \vee C \vdash B \vee C$
4. *Deduction Theorem.* If $\Gamma, A \vdash B$, then $\Gamma \vdash A \to B$
5. *Proof by contradiction.* $\Gamma, \neg A \vdash \bot$ iff $\Gamma \vdash A$
6. *Post's Theorem.* (Also called "tautology theorem", or even "completeness of Propositional Calculus theorem")

   If $\models_{\text{taut}} A$, then $\vdash A$.

   Also: If $\Gamma \models_{\text{taut}} A$, then $\Gamma \vdash A$.
7. *Proof by cases.* $A \to B, C \to D \vdash A \vee C \to B \vee D$

   Also the special case: $A \to B, C \to B \vdash A \vee C \to B$

Translations

$(\exists \mathbf{x})A$ translates to $\neg(\forall \mathbf{x})\neg A$

$(\forall \mathbf{x})_B A$ translates to $(\forall \mathbf{x})(B \to A)$

$(\exists \mathbf{x})_B A$ translates to $(\exists \mathbf{x})(B \wedge A)$

Useful facts from Predicate Calculus (proved in class—you may use them without proof):

We ***know*** that SL and WL are **derived rules** useful in equational proofs within predicate calculus.

► More "rules" and (meta)theorems.

(i) *Dummy renaming.*

If $\mathbf{z}$ does not occur in $(\forall \mathbf{x})A$ as either free or bound, then $\vdash (\forall \mathbf{x})A \equiv (\forall \mathbf{z})\big(A[\mathbf{x} := \mathbf{z}]\big)$

If $\mathbf{z}$ does not occur in $(\exists \mathbf{x})A$ as either free or bound, then $\vdash (\exists \mathbf{x})A \equiv (\exists \mathbf{z})\big(A[\mathbf{x} := \mathbf{z}]\big)$

(ii) *$\forall$ over $\circ$ distribution, where "$\circ$" is "$\vee$" or "$\rightarrow$".*

$\vdash A \circ (\forall \mathbf{x})B \equiv (\forall \mathbf{x})(A \circ B)$, **provided x** is not free in A

$\exists$ over $\wedge$ distribution

$\vdash A \wedge (\exists \mathbf{x})B \equiv (\exists \mathbf{x})(A \wedge B)$, **provided x** is not free in A

(iii) *$\forall$ over $\wedge$ distribution.*

$$\vdash (\forall \mathbf{x})A \wedge (\forall \mathbf{x})B \equiv (\forall \mathbf{x})(A \wedge B)$$

$\exists$ over $\vee$ distribution.

$$\vdash (\exists \mathbf{x})A \vee (\exists \mathbf{x})B \equiv (\exists \mathbf{x})(A \vee B)$$

(iv) *$\forall$ commutativity (symmetry).*

$$\vdash (\forall \mathbf{x})(\forall \mathbf{y})A \equiv (\forall \mathbf{y})(\forall \mathbf{x})A$$

$\exists$ commutativity (symmetry).

$$\vdash (\exists \mathbf{x})(\exists \mathbf{y})A \equiv (\exists \mathbf{y})(\exists \mathbf{x})A$$

(v) *Specialisation. "Spec"* $(\forall \mathbf{x})A \vdash A[\mathbf{x} := t]$, for any term t.

Dual of Specialisation. $A[\mathbf{x} := t] \vdash (\exists \mathbf{x})A$, for any term t.

(vi) *Generalisation. "Gen"* If $\Gamma \vdash A$ and if, moreover, the formulae in Γ have **no free x occurrences**, then also $\Gamma \vdash (\forall \mathbf{x})A$.

(vii) *$\forall$ Monotonicity.* If $\Gamma \vdash A \rightarrow B$ so that the formulae in Γ have **no free x occurrences**, then we can infer

$$\Gamma \vdash (\forall \mathbf{x})A \rightarrow (\forall \mathbf{x})B$$

(viii) $\forall$ *Introduction; a special case of* $\forall$ *Monotonicity.* If $\Gamma \vdash A \to B$ so that neither the formulae in Γ nor A have **any free x occurrences**, then we can infer

$$\Gamma \vdash A \to (\forall \mathbf{x})B$$

(ix) Finally, the *Auxiliary Variable ("witness") Metatheorem.* If $\Gamma \vdash (\exists \mathbf{x})A$, and if $\mathbf{y}$ is a variable that ***does not*** occur as either free or bound variable in any of $(\exists \mathbf{x})A$ or B or the formulae of Γ, then

$$\Gamma, A[\mathbf{x} := \mathbf{y}] \vdash B \text{ implies } \Gamma \vdash B$$

Semantics facts

Propositional Calculus	Predicate Calculus
(Boolean Soundness) $\vdash A$ implies $\models_{\text{taut}} A$	$\vdash A$ does **NOT** imply $\models_{\text{taut}} A$
(Post) $\models_{\text{taut}} A$ implies $\vdash A$	However, (Post) $\models_{\text{taut}} A$ implies $\vdash A$
	(Pred. Calc. Soundness) $\vdash A$ implies $\models A$

CAUTION! The above facts/tools are only a fraction of what we have covered in class. They are ***very important and very useful***, and that is why they are listed for your reference here.

You can also use *without proof* ***ALL*** the things we have covered (such as the absolute theorems known as "$\exists$-definition", "de Morgan's laws", etc.).

But these —the unlisted ones— are up to you to remember and to correctly state!

Whenever in doubt of whether or not a "tool" you are about to use was indeed covered in class, ***prove*** *the validity/fitness of the tool before using it!*