Lassonde Faculty of Engineering **EECS**

MATH1090B. Problem Set No1 Posted: Sept. 15, 2018

Due: Oct. 3, 2018, by 2:30pm; in the course assignment box.



It is worth remembering (quoted from the course outline):

The homework must be each individual's own work. While consultations with the instructor, tutor, and among students, are part of the learning process and are encouraged, nevertheless, at the end of all this consultation each student will have to produce an individual report rather than a copy (full or partial) of somebody else's report.

The concept of "late assignments" does not exist in this course.



- 1. (5 MARKS) Take as given [done in class] that every wff contains an atomic formula as a substring—as we say, it has an atomic subformula. Then prove by analysing formula-calculations or, preferably, using the recursive definition of wff structure, that $(p \rightarrow)$ is not a wff.
- 2. (6 MARKS) Prove via formula calculations or using the recursive definition of wff that
 - No wff can end with the symbol \equiv .
 - No wff can contain the substring $\equiv \vee$.
- **3.** (1 MARK) Prove that $(r \wedge ((\neg (p \equiv q)) \wedge p)) \rightarrow (r' \rightarrow \bot)$ is a wff.
- 4. (6 MARKS) Recall that a schema is a tautology iff all its instances are tautologies. Thus,

A schema is not a tautology iff it has an instance that is not.



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Which of the following six schemata are tautologies? Show the whole process that led to your answers, *including truth tables or* equivalent short cuts, *and words of explanation*.

I note that in the six sub-questions below I am *not* always using all the formally necessary brackets. It is your task to correctly insert any missing brackets before you tackle the question for each formula.

- $A \equiv B \equiv A \wedge B \vee \neg A \wedge \neg B$
- $A \equiv B \equiv A \equiv \bot \equiv B \equiv \bot$
- $A \lor B \to A \land B$
- $A \to B \equiv \neg B \to \neg A$
- $(A \equiv B) \rightarrow A \land B$
- $A \equiv B \equiv (A \rightarrow B) \land (B \rightarrow A)$
- **5.** (5 MARKS) By using truth tables, or using related shortcuts, examine whether or not the following *tautological implications* are correct.
- In order to show that a tautological implication that involves *meta*-variables for formulae —i.e., it is a schema— is *incorrect* you *must* consider an *instance* (i.e., a special case with specific formulae) that *is* incorrect.



Show the whole process that led to each of your answers.

- $p \models_{\text{taut}} p \land q$
- $\top \models_{\text{taut}} A$
- $\top \models_{\text{taut}} A \to A$
- $\neg A \land A \models_{\text{taut}} B$
- $A, B \to A \models_{\text{taut}} B$
- **6.** (6 MARKS) Write down the *result* of the following substitutions, whenever the requested substitution makes sense. The result must not contain the string " $[\mathbf{p} := \dots]$ ".

Whenever a requested substitution does $\underline{\text{not}}$ make sense, explain $\underline{\text{exactly}}$ why it does not.

Show the whole process that led to each of your answers in each case.

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Remember the priorities of the various connectives as well as that of the meta-expression " $[\mathbf{p}:=\ldots]$ "! The following formulae have not been written with all the formally required brackets.



- $p \to \top[p := \mathbf{t}]$
- $\bullet \ (p \to r)[q' := \bot]$
- $p \lor q \land r[r := A]$ (where A is some formula)
- $p \lor (q \to p)[p := r]$
- $(p \lor q)[p := \mathbf{t} \to \mathbf{f}]$
- $(p \lor q)[(p \lor q) := r]$