## Department of EECS

MATH1090 B. Problem Set No1 Posted: Sept. 19, 2021

Due: Oct. 5, 2021, by 2:00pm; in eClass.

Q: How do I submit?

**A**:

- (1) Submission must be ONLY ONE file
- (2) Accepted File Types: PDF, RTF, MS WORD, ZIP
- (3) Deadline is strict, electronically limited.
- (4) MAXIMUM file size = 10MB
- It is worth remembering (from the course outline):

The homework must be each individual's <u>own work</u>. While consultations with the <u>instructor</u>, tutor, and <u>among students</u>, are part of the <u>learning process</u> and are encouraged, nevertheless, at the end of all this consultation each student will have to produce an <u>individual report</u> rather than a copy (full or partial) of somebody else's report.

The concept of "late assignments" does not exist in this course.



1. (3 MARKS) Prove that no wff can be  $\lambda$ .

*Hint.* Analyse formula-constructions, or use induction on formulas.

Page 1 G. Tourlakis

2. (3 MARKS) Prove that the complexity of a well-formed-formula equals the number of its right brackets.

Hint. Analyse formula-constructions, or use induction on formulas.

- **3.** (3 MARKS) Prove that (**p**) is *not* a wff. *Hint.* One way is to analyse formula-constructions.
- **4.** (1 MARK) Prove that  $((\neg(p \to r)) \equiv (r \to p))$  is a wff.
- **5.** (6 MARKS) Recall that a schema is a tautology iff *all* its *instances* are tautologies.

Which of the following six schemata are tautologies? Show the whole process that led to your answers, including truth tables or equivalent short cuts, and words of explanation.

I note that in the six sub-questions below I am <u>not</u> using <u>all</u> the <u>formally</u> necessary brackets.



Therefore be mindful of connective priorities and associativities!



- $A \to B \to (A \to \bot) \lor B$
- $A \equiv B \rightarrow (A \rightarrow \bot) \lor B$
- $(A \equiv B) \rightarrow A \land B$
- $A \to B \to \neg B \to \neg A$
- $(\neg A) \land B \equiv A \to B$
- $A \vee B \rightarrow A \rightarrow B$
- **6.** (3 MARKS) Prove that if we have  $\top \models_{\text{taut}} A$ , then we also have  $B \models_{\text{taut}} A$  for any B.
- 7. (6 MARKS) By using truth tables, or using related shortcuts, examine whether or not the following tautological implications are correct.
- In order to show that a tautological implication that involves *meta*-variables for formulas —i.e., it is a schema— is *incorrect* you *must* consider a special case that *is* incorrect (since some other special cases might work).



Page 2 G. Tourlakis

Show the whole process that led to each of your answers.

- $p \land \neg p \models_{\text{taut}} \bot$
- $p \lor q \land r \to r'' \models_{\text{taut}} \top$
- $p \models_{\text{taut}} p \vee B$
- $A, A \rightarrow B \models_{\text{taut}} B$
- $A \equiv B \models_{\text{taut}} A \wedge B$
- $A \wedge B \models_{\text{taut}} A \equiv B$
- 8. (6 MARKS) Write down the most simplified result of the following substitutions, whenever the requested substitution makes sense. Whenever a requested substitution does <u>not</u> make sense, explain <u>exactly</u> why it does not.

Show the whole process that led to each of your answers in each case.

Remember the priorities of the various connectives as well as that of the meta-expression " $[\mathbf{p} := \dots]$ "! The following formulas <u>have not</u> been written with all the formally required brackets.



- $(q \to p)[q := r]$
- $(q \rightarrow p)[r := r'']$
- $\bullet \ p \to \top[p := \bot]$
- $p \to \top[p := \mathbf{f}]$
- $(\bot \to r \to q)[\bot := p]$
- $\bullet \ p \vee (q \wedge r)[p := r]$