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30 March 2009

ICDE



for XML in DB2

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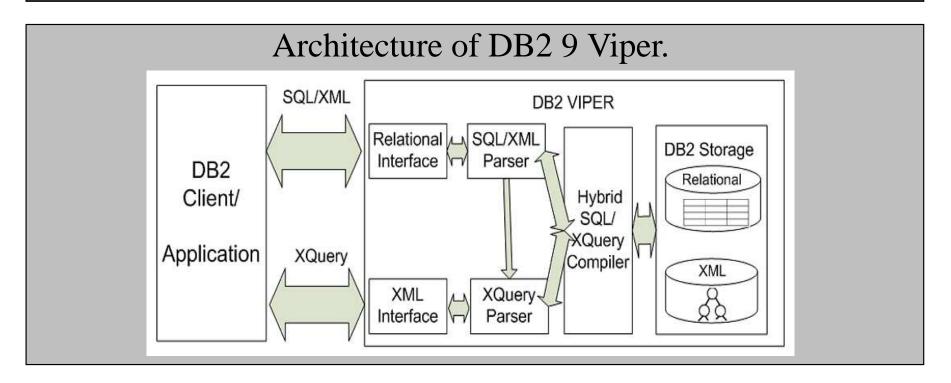
In today's big, commercial database systems, relational and XML stand side by side.

Much work has been done

- to handle XML efficiently,
 - native XML storage
 - XML indexes
- and to *integrate* relational and XML data
 - SQL can query XML data
 - XQuery can query relational data



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Logical integration of XML in relational

- XML is a data type for columns
 - each value is an XML document
 - an XML column can be contrained by XML Schemas.
- table function XMLTable extends SQL
 - part of the SQL standards (SQL/XML)



In today's big, commercial database systems, relational and XML stand side by side.

Challenge

- SQL and relational database systems have nearly 40 years of optimization technology.
- Queries over XML (and hybrid) can be quite inefficient.
 - many fewer robust optimization techniques
 - queries are inherently "more complex"

Can we adapt successful techniques for relational queries for XML (and hybrid)?

Optimization

Materialized Views

A *materialized view* is simply a view that is evaluated and stored as a table on disk. It is maintained and updated by the database system.

- This can speed up the evaluation of other queries.
- But updates to the tables in the view's definition will (usually) require updates to the view's table.

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The technology of materialized views

- is quite well studied, and
- is effectively implemented and used in relational database systems.

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Q. Can we do the same for views that use XMLTable to optimize SQL/XML queries?

```
SELECT x.title, x.sections
FROM books,
XMLTable ('$book/book/chapter'
PASSING bookdoc AS "book"
COLUMNS
    "title" VARCHAR (60)
                          PATH 'title',
                          "sections" XML PATH 'section'
) AS x;
```

<?xml version="1.0" ...?> <book> <chapter> <title>Introduction</title> <section>Goals</section> </chapter> <chapter> <title>Related</title> <section>Views</section> <section>XML</section> </chapter> <chapter> <title>Solutions</title> </chapter> </book>

SELECT x.title, x.sections
FROM books,
XMLTable ('\$book/book/chapter'
PASSING bookdoc AS "book"
COLUMNS
 "title" VARCHAR (60)
 PATH 'title',
 "sections" XML PATH 'section'
) AS x;

• row generator: Generates rows based on documents in the XML column. <?xml version="1.0" ...?> <book> <chapter> <title>Introduction</title> <section>Goals</section> </chapter> <chapter> <title>Related</title> <section>Views</section> <section>XML</section> </chapter> <chapter> <title>Solutions</title> </chapter> </book>

• row generator: Generates rows based on documents in the XML column.

• navigtors: Populate the columns' values for those rows.

<?xml version="1.0" ...?> <book> <chapter> <title>Introduction</title> <section>Goals</section> </chapter> <chapter> <title>Related</title> <section>Views</section> <section>XML</section> </chapter> <chapter> <title>Solutions</title> </chapter> </book>

SELECT x.title, x.sections			xml version="1.0"?>
FROM books,		<b< td=""><td>ook></td></b<>	ook>
XMLTable ('\$book/book/chapter'			<chapter></chapter>
PASSING bookdoc AS "book"			<title>Introduction</title>
COLUMNS			<section>Goals</section>
"title" VARCHAR (60)			
PATH 'title',			<chapter></chapter>
"sections" XML PATH 'section'			<title>Related</title>
) AS x;			<pre><section>Views</section></pre>
			<section>XML</section>
			chapter>
title	sections		hapter>
Introduction	<pre><section>Goals</section></pre>	m>	<title>Solutions</title>
		/11/	chapter>
	<section>Views</section>		ok>
Related	<section>XML</section>		

Solutions

null

The Challenges

A matching and compensation framework is needed.

- *matching:* discovering whether the view is equivalent to, or contains, the query.
- *compensation:* determine further restrictions—navigation steps and predicates—which, when applied to the view, is equivalent to the query.

• The Challenges

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related work

• XPath rewrites were investigated in

[Balmin, Özcan, Beyer, Cochrane, & Pirahesh, VLDB 2004]

[Xu & Özsoyoglu, VLDB 2005]

• and for SQL/XML in

[Krishnaprasad, Liu, Manikutty, Warner, Arora, & Kotsovolos, VLDB 2004]

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Considerations

- The row generator of the query and view *must* be the same.
- The navigators must
 - 1. be type convertible,
 - 2. satisfy the prefix property (be less restrictive), and
 - 3. satisfy compensation locality.

• The Challenges

A matching and compensation framework is needed.

- *matching:* discovering whether the view is equivalent to, or contains, the query.
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A step is less restrictive than another if

- the steps have the same axis navigators,
- the step's test is less restrictive than the other's, and
- the step's predicate implies the other's.

Implementation

The existing matching and compensation framework in DB2 9.5 (VIPER 2) was extended to handle new types of matching using XMLTable.

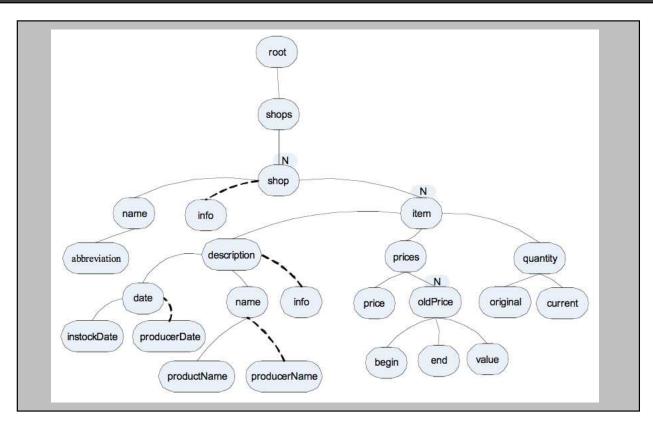
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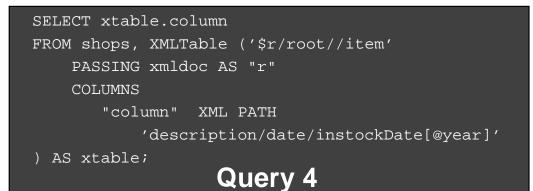
- Developed a subsumption algorithm.
 - Given a "subsumer" tree of size m and a "subsumee" tree of size n, takes mn steps.
- Relaxed strict *locality*.
- Integrated within DB2's optimizer's materialized view mechanism.
- Addressed current "impedance" issues.

Testing Environment

- A code branch of IBM DB2 9.5 was developed, implementing the above.
- 13 XMLTable queries were devised based on the Shops schema to test the different subsumption patterns.
- A 10MB and a 100MB Shops databases were synthetically generated.



Query Suite



Query Suite

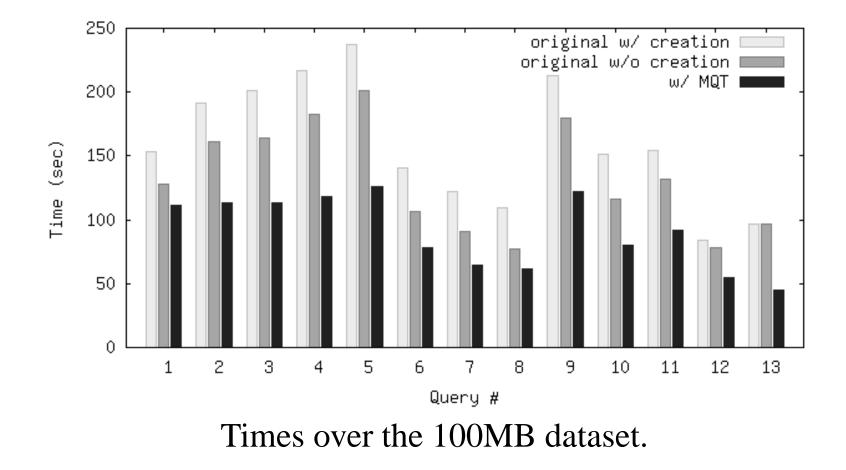
#	Row Generator	Query Navigator	
#	MQT Navigator	Compensation	
1	/root/shops/shop/item	prices/price	
	prices	prices/price	
2	/root//item	prices/price	
2	prices	prices/price	
3	/root//item	description/date/instockDate	
	description	description/date/instockDate	
4	/root//item	description/date/instockDate[@year]	
4	description/date	date/instockDate[@year]	
5	/root//item	description/date/instockDate[@year and @month and @day]	
5	description/date	date/instockDate[@year and @month and @day]	
6	/root//item	description[info]/date/instockDate[@year>2006]	
0	description	<pre>description[info]/date/instockDate[@year>2006]</pre>	

Query Suite

#	Row Generator	Query Navigator	
#	MQT Navigator	Compensation	
7	/root//item	<pre>description[contains(info,"car")]/name</pre>	
	description	<pre>description[contains(info,"car")]/name</pre>	
8	/root//item	description/date/productionDate[@year>2005]//instockDate	
0	description	<pre>description/date/productionDate[@year>2005]//instockDate</pre>	
9	/root//item	description[date]/name/productName	
	description[date]	description/name/productName	
10	/root//item	description[info and date]/name/productName	
10	description[date]	description[info and date]/name/productName	
11	/root/*/*/item	description/date	
	description	description/date	
12	/root//*//shop	name/abbreviation	
12	name/*	abbreviation	

Experimental Results

& Lessons Learned



III. Conclusions & Future Work

conclusions

- This work offers a strong evidence that materialized views can be highly effective for SQL/XML queries.
- Is the first work, to the best of our knowledge, that considers use of materialized views to improve SQL/XML performance for XML data.
- While our present matching and compensation is quite simple, it seems applicable for a wide range of SQL/XML queries.

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future work

- 1. Extend the matching and compensation framework for a wider class of XML/SQL queries.
- 2. Move the framework into the cost-based optimizer.
- 3. Make the framework applicable to XQuery.