

Database Management Systems

York University

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EECS-3421 vs EECS-4411

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Ha! No way.

In this class, we focus on how to *build* a database system. In EECS-3421, we focused on what functionality a database system provides, and how to *use* it.

Data Independence

Do not need to know how...

- a *compiler* works to write a program.
- an *operating system* is built to use one.
- a *car* works to drive one.
- a *database system* is built to use it.

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- **physical data independence:** how the data is *logically* organized is independent of how it is *physically* organized. (There is also *logical data independence*...)
- **Codd's law**: Can only access and update the database via the "query language" (SQL).
- SQL is a *declarative* language.

How to build a Database System?

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In this class, we're going to build our own system!

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- What are our *design criteria*?
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- What are our *design choices*? Our *design constraints*?
 - How will the available technology affect our design (*architecture*)?

E.g., Main memory technologies (like CMOS) are volatile.



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Goals:

- permanence
- fast, random access
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Design questions:

- What devices / technology do we use?
- What data-structures do we use? How do we access given pieces of data quickly?



How to evaluate (SQL) queries efficiently? Need a

- query parser
- plan generator / query optimizer *Turns a valid SQL query into a "program" that answers the query.*
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Problems:

- SQL is reasonably complex.
- Not all (equivalent) queries are equal. Some queries / query plans will evaluate inherently must faster.

Big issue:

• How to "pick", or design, a good query plan for a query?

A "Complex" Query

Supplier **S**: A (name), C (city) Retailer **R**: B (name), C (city)

Query: Which supplier has a location in every city of a retailer? Show such supplier (A) / retailer (B) pairs.

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A "Complex" Query in SQL

select A, B from R, S except select A, B from (select S.A, R.B, R.C from R, S except select S.A, R.B, R.C from R, S where R.C = S.C) as Z;

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Any problems?

A "Complex" Query Better?

```
select A, B
        from R, S
        where \mathbf{R}.\mathbf{C} = \mathbf{S}.\mathbf{C}
except
select A, B from (
        select S.A, R_1.B, R_2.C
                 from R as \mathbf{R}_1, R as \mathbf{R}_2, S
                 where \mathbf{R}_1 \cdot \mathbf{C} = \mathbf{S} \cdot \mathbf{C} and \mathbf{R}_1 \cdot \mathbf{B} = \mathbf{R}_2 \cdot \mathbf{B}
        except
        select S.A, R.B, R.C from R, S
                 where \mathbf{R}.\mathbf{C} = \mathbf{S}.\mathbf{C}
) as Z;
```

A "Complex" Query cleaned up

with

J (A, B, C) as (select S.A, R.B, R.C from **R**, **S** where $\mathbf{R}.\mathbf{C} = \mathbf{S}.\mathbf{C}$) select distinct A, B from J except select J.A, J.B from J, R where $\mathbf{J}.\mathbf{B} = \mathbf{R}.\mathbf{B}$ and (**J**.A, **J**.B, **R**.C) not in (select A, B, C from J);

A "Complex" Query via COUNT

select J.A, J.B from (select S.A, R.B, count(*) as Cs from R, S where R.C = S.Cgroup by S.A, R.B) as J, (select B, count(*) as Cs from R group by B) as K where J.B = K.B and J.Cs = K.Cs;

The Query Optimizer

Rewrite

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Cost-based

- a. Determine a "best" over-all query tree.
- b. Pick the best method for each operator in the query tree.
 - 1) Pick the best *access path* for each table involved.
 - 2) Assign the "best" algorithm to each operator

 $(\bowtie, \pi, \sigma, \ldots).$

c. Do **a.** & **b.** simultaneously!



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Properties: • A • C • I • D

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Properties:

- Atomicity
- Consistency
- Isolation
- **D**urability

Building a Database System

Anything we miss?

- host language support e.g., JDBC
- data definition language (DDL)
 e.g., CREATE TABLE ...
- administrative functions (for DBA's) & security e.g., GRANT ...

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• . . .

What pieces / modules do we need to implement all this? What's our architecture? Need a

- query optimizer
- transaction manager
 - lock manager for concurrency control
- crash recovery mechanism

)..



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- Cannot be a *good* DB Administrator *without* understanding how the system works.
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- Lots of places are building database-like systems. *Can reuse the techniques and technologies from RDBMSs.*